

Exercise Sheet 3: Specification and Verification with Higher-Order Logic (Summer Term 2014)

Please prepare the marked tasks for the exercise on Wednesday, May 21, 2014
Submit your solutions to the hand-in tasks before Wednesday, May 28, 2014

Exercise 1 Methods and Rules in Isabelle/HOL

a) (Prepare!) Apply the rule

$$\llbracket (\lambda a. \lambda b. (a, b) \in r^*; \bigwedge x. ?P x x; \bigwedge x y z. \llbracket (x, y) \in r^*; ?P x y; (y, z) \in r \rrbracket \implies ?P x z \rrbracket \implies ?P a b$$

with the method `erule` to the following subgoal by hand (i.e. on paper):

$$(i, j) \in s^* \implies 0 \leq (\text{dist } i j)$$

Hint: Don't be distracted by unknown function names; you don't have to know anything about their meaning. Just apply the rule syntactically.

b) In this exercise we want to practice the use of different methods (like `rule`, `erule` or `frule`) to prove properties in propositional and predicate logic. You should only use the methods `rule`, `erule`, `frule`, `drule`, the respective `_tac` methods and `assumption`. Do **not** use other methods like `simp`. You should only use the rules of the first exercise sheet, together with the following additional rules: `conjE`, `impE`, `iffI`, `iffE`, and `classical`.

Hint: You can write “`thm classical`” in Isabelle/HOL to see the concrete definition of the rule.

Prove the following theorems:

1. $A \wedge B \longrightarrow B \wedge A$
2. $(A \vee A) = (A \wedge A)$
3. $A \longrightarrow B \longrightarrow A$
4. $(A \longrightarrow (B \longrightarrow C)) \longrightarrow ((A \longrightarrow B) \longrightarrow (A \longrightarrow C))$
5. $(\neg A \longrightarrow \neg B) \longrightarrow (B \longrightarrow A)$
6. $\neg\neg A \longrightarrow A$
7. $A \vee \neg A$
8. $(\exists x. \forall y. P x y) \longrightarrow (\forall y. \exists x. P x y)$
9. $((\forall x. P x) \wedge (\forall x. Q x)) = (\forall x. (P x \wedge Q x))$
10. $((\exists x. P x) \vee (\exists x. Q x)) = (\exists x. (P x \vee Q x))$
11. $(\neg(\forall x. P x)) = (\exists x. \neg P x)$

Exercise 2 Rewriting and Simplification

In this exercise we want to do proofs by just rewriting. Please download the file `Sheet3_Rewrite.thy` from the website. You are only allowed to use the lemmas defined in this file and you are only allowed to use them with the `subst` method. As the only exception you are allowed to use `(rule TrueI)` to finish a subgoal.

Prove the following theorems:

- a) $\text{length } [a] = 1$
- b) $(A \wedge B \wedge C) = (B \wedge A \wedge C)$
- c) $(a * (\text{Suc } (\text{Suc } (\text{Suc } (\text{Suc } 0)))) \text{ div } (\text{Suc } (\text{Suc } (\text{Suc } (\text{Suc } 0)))) = a$
- d) $(\text{Suc } x) * (\text{Suc } x) = \text{Suc } (x*x + 2*x)$

Exercise 3 Prime numbers (Hand in!)

Please download the file `Sheet3_Primes.thy` from the website. This file contains an unfinished proof for a theorem, which states that there is an infinite number of primes. Your task is to finish this proof by formalizing the following informal proof:

Lemma: For every number greater or equal to 2 there exists some prime which divides the number.

Proof: By induction over n .

Induction Hypothesis: For every number k between 2 and $n - 1$ there exists some prime which divides k .

Induction Step (show the statement for n using the Induction Hypothesis): If n is prime the step is trivial. If n is not prime, then there exists a number k with $2 \leq k < n$ which divides n (by the definition of prime number). From the induction hypothesis we get a prime number k' which divides k . By transitivity of “divides” k' also divides n . So with k' we have a prime number dividing n .

Theorem: There are infinitely many primes.

Proof: Suppose for the sake of contradiction that the set S of primes is finite. Then $\prod S$ is well defined. Let $P = 1 + \prod S$. Then there exists some prime q which divides P . Because $q \in S$, q also divides $\prod S$. If a number divides two numbers n and m , then it also divides the difference $m - n$. Therefore q divides the difference $P - \prod S$, which is 1. Only 1 divides 1, so $q = 1$, but it also is a prime number. This is a contradiction and the proof is complete.

Hints:

- $\prod S$ denotes the product of all numbers in the set S , for example $\prod\{2, 3, 5\} = 2 \cdot 3 \cdot 5 = 30$.
- In Isabelle $n \text{ dvd } m$ denotes the fact that n divides m , for example $3 \text{ dvd } 12$.
- The induction rule used for the first lemma is called `full_nat_induct`. The difference to the usual induction over natural numbers is, that one can assume, that the hypothesis holds for all numbers smaller than n , whereas in the usual induction rule it can only be assumed for the direct predecessor.
- The following theorems might be helpful: `dvd_diff_nat`, `dvd_setprod`, `dvd_trans`.
- `Collect P` is the set of all elements for which P holds true.