

Exercise Sheet 10: Specification and Verification with Higher-Order Logic (Summer Term 2012)

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Exercise 1 Hoare Calculus

Consider the following while program `prog`, which calculates the n th Fibonacci number in linear time:

```
f0 := 0;
f1 := 1;
i := 0;

WHILE i < n INV ? DO
  f1 := f0 + f1;
  f0 := f1 - f0;
  i := i + 1
END;

result := f0
```

As a postcondition we would like to have that `result` is `fib n`, where `fib` is specified as usual:

$$\text{fib } n = \begin{cases} 0 & n \leq 0 \\ 1 & n = 1 \\ \text{fib } (n - 1) + \text{fib } (n - 2) & \text{otherwise} \end{cases}$$

a) (Prepare!) Find a suitable loop invariant and prove the following Hoare-triple *on paper*.

$$\{n \geq 0\} \text{ prog } \{\text{result} = \text{fib } n\}$$

- b) Translate the proof into Isabelle/HOL, using the `Hoare.thy` from the website. Apply all Hoare rules manually, like done in the `sqrProp` lemma.
- c) Optimize your proof by using the custom `wphoare` method, like done in the `splitcorrect` lemmas.
- d) The `splitcorrect` lemma does not prove that `split` does not change the content of the array.

At the end of the theory file, we have specified a `content` function for arrays in a state. Using this function, we would like to prove the following Hoare-triple:

$$\{\text{splitRequires} \wedge \text{content myarr} = C\} \text{ splitProg } \{\text{content myarr} = C\}$$

Find suitable invariants for the loops and prove the triple using `wphoare`.